

**AMENDMENT TO THE CLAIMS**

This listing of claims will replace all prior versions of claims in the application.

**Listing of Claims:**

1. (currently amended) A system for managing coherent data access through multiple nodes, comprising:

a first data processing system forming a first node, in which the first data processing system includes a first processor, a first bridge, a first interface and a first memory that is local to the first node, wherein the first node supports packet traffic for transfer of packets, coherent traffic to access the first memory and non-coherent traffic to communicate with input/output (I/O) circuitry, in which the first memory stores first cacheable data having global coherency; and

a second data processing system forming a second node that also supports packet traffic, coherent traffic and non-coherent traffic, in which the second data processing system includes a second processor, a second bridge, a second interface and a second memory that is local to the second node—and, the second memory storing second cacheable data also having global coherency, the first and second interfaces coupling the first node to the second node via a bi-directional bus for transfer of packet traffic, coherent traffic and non-coherent traffic between the first and second nodes, wherein when the second node receives a data packet from an external source that is to access a cache coherent fabric of the first memory, the second bridge identifies the first memory to be as located in-a-remote the first node and converts the data packet to an uncacheable data access request to the first node instead of performing a cache coherent memory access operation to access the first memory from the second node, so that the access to the first memory does not access a coherent fabric of the second memory in the second node, and when the first bridge receives the uncacheable data access request, the first bridge identifies the uncacheable data access request as a local access request to the first memory in the first node and processes the uncacheable data access request from the second node now as a coherent data access request in the first node to access the coherent fabric of the first memory in the first node.

2. (canceled)
3. (previously presented) The system of claim 1, wherein the uncacheable data access request associated with the packet from the external source is a store or a write request to access the first memory.

- 4-5. (canceled)

6. (previously presented) The system of claim 1, wherein the uncacheable data access request by the second bridge follows a producer-consumer protocol.
7. (currently amended) The system of claim 6, wherein a payload and a flag are written to a memory location ~~in-a-home~~ the first node when following the producer-consumer protocol.

- 8-9. (canceled)

10. (currently amended) A method for managing coherent data access through multiple nodes, comprising:

establishing a cacheable coherent memory space in a first memory of a first data processing system that forms a first node, the first memory being local to the first node and to store first cacheable data having global coherency, wherein the first node supports packet traffic for transfer of packets, coherent traffic to access the first memory and non-coherent traffic to communicate with input/output (I/O) circuitry, and in which the first data processing system also includes a first processor, a first bridge and a first interface;

receiving at a second node a data packet from an external source to access a cache coherent fabric of the first memory in the first node, wherein the second node also supports packet traffic, coherent traffic and non-coherent traffic and is formed of a second data processing system that includes a second processor, a second bridge, a second interface and a second memory that is local to the second node-and, the second memory storing second cacheable data also having global coherency, the first and second

interfaces coupling the first node to the second node via a bi-directional bus for transfer of packet traffic, coherent traffic and non-coherent traffic between the first and second nodes;

identifying in the second node that the first memory—that is to be accessed is located in-a-remote the first node;

converting the data packet in the second bridge to an uncacheable data access request to access the first memory in the first node instead of performing a cache coherent memory access operation to access the first memory from the second node, so that the access to the first memory does not access a coherent fabric of the second memory in the second node;

transferring by the second bridge the uncacheable data access request to the first node;

receiving by the first bridge the uncacheable data access request from the second node through coupling between the first and second interfaces;

identifying the uncacheable data access request to the first memory as a local access in the first node;

processing the uncacheable data access request from the second node now as a coherent data access to access the coherent fabric of the first memory in the first node; and

accessing the coherent fabric of the first memory in the first node in response to receiving the packet from the external source, but without accessing the coherent fabric of the second memory in the second node.

11. (canceled)

12. (previously presented) The method of claim 10, wherein the uncacheable data access request associated with the packet from the external source is a store or a write request to access the first memory.

13-14. (canceled)

15. (previously presented) The method of claim 10, wherein the uncacheable data access request by the second bridge follows a producer-consumer protocol.
16. (currently amended) The method of claim 15, wherein a payload and a flag are written to a memory location ~~in-a-home~~ the first node when following the producer-consumer protocol.

17-18. (canceled)